

A Novel Linear Algebra Method for the Determination of Periodic Steady States of Nonlinear Oscillators

Abstract—Periodic steady-state (PSS) analysis of nonlinear oscillators has always been a challenging task in circuit simulation. We present a new way that uses numerical linear algebra to identify the PSS(s) of nonlinear circuits. The method works for both autonomous and excited systems. Using the harmonic balancing method, the solution of a nonlinear circuit can be represented by a system of multivariate polynomials. Then, a Macaulay matrix based root-finder is used to compute the Fourier series coefficients. The method avoids the difficult initial guess problem of existing numerical approaches. Numerical examples show the accuracy and feasibility over existing methods.

Keywords—Steady-state analysis, autonomous oscillator, nonlinear circuit simulation, Macaulay matrix

I. INTRODUCTION

Accurate simulation of the behaviors of microwave and radio-frequency (RF) circuits has become a key issue in mixed-signal circuit design due to the exponential growth of the complexity in the past few years. Periodic steady-state (PSS) analysis plays a critical role, especially in the simulation of RF communication circuits [1].

Two specialized PSS algorithms, namely, shooting Newton (SN) and harmonic balance (HB) [2]–[5] are deployed in traditional SPICE simulators for decades. They are well developed for the PSS analysis of excited systems, i.e., systems with certain periodic inputs. In such systems, the period of the PSS is always forced to be the same as that of the input. SN will then iteratively shoot for aligning the initial and final values over the given period, while HB numerically solves for the Fourier series coefficients with respect to the known period. However, it is difficult for traditional SN and HB to work under an unknown oscillating frequency. Therefore, they cannot readily be used to determine the PSS of autonomous systems (i.e., systems without external inputs) such as oscillators in RF communication circuits.

Several modifications have been made for SN and HB to adapt to autonomous systems by treating the unknown oscillation frequency as a new state variable [6]–[9]. These methods still require solving systems of nonlinear equations via the Newton-Raphson method, which highly depends on the initial guess. There might be multiple solutions to the nonlinear equations where only one is stable. Therefore it is always hard to choose an initial value of the frequency in the right basin of attraction, especially for RF oscillators.

To avoid dependence on the initial guess when solving nonlinear systems numerically, other methods, such as Volterra series [10], [11], have been proposed. The oscillation frequency is then obtained by solving a univariate equation which

consists of Volterra transfer functions. However, Volterra series approximation might be inaccurate, especially for highly nonlinear oscillators. Moreover, Volterra transfer functions suffer from singularity problems when applied to descriptor systems. At the same time, algebraic geometry has been used to solve polynomial systems without any initial guesses. For example in [12], a polynomial system of harmonic balancing equations is solved using a Groebner basis. Nonetheless, although no initial guesses are needed in the method, Groebner bases can only deal with integer coefficients. In addition, Groebner bases typically generate polynomials with very large coefficients (usually hundreds of digits), which is generally inconvenient to use in engineering.

On the other hand, a new numerical way of finding all the roots of a multivariate polynomial system has been recently introduced in [13], [14]. This method uses only numerical linear algebra and works with arbitrary real coefficients while finding all the roots by computing the null space of the Macaulay matrix. The main computational tools in this method are either the singular value decomposition (SVD) or the rank-revealing QR decomposition, for which numerically backward stable implementations exist [15].

In this paper, a novel linear algebraic method is proposed for the PSS analysis of nonlinear oscillators. First, a general nonlinear system is transformed into an equivalent polynomial differential algebraic equation (DAE) system. Then, a polynomial system is constructed from the DAE system via harmonic balancing equations. After that, the Macaulay matrix based root-finder is used to solve for the Fourier series coefficients from the polynomial system. The proposed method completely avoids initial guesses as in iterative numerical solvers such as SN and HB, and works perfectly with arbitrary system coefficients. Moreover, it is able to find all solutions as long as the PSS of the system exists.

The paper is organized as follows. Section II reviews the background preliminaries of the polynomial DAE system. Then, our numerical linear algebra method is proposed in Section III. Numerical examples are given in Section IV. Several notable remarks are presented in Section V. Section VI concludes this paper.

II. BACKGROUND AND PRELIMINARIES

A. Systems of polynomial differential algebraic equations

First, we start with a general nonlinear DAE system

$$F(x, \dot{x}, \ddot{x}, \dots, u) = \mathbf{0}, \quad (1)$$

where $x \in \mathbb{R}^n$, $u \in \mathbb{R}^l$ are the state vector and input vector, respectively. $F(\cdot)$ is an arbitrary nonlinear vector function. In practical circuit modeling, $F(\cdot)$ usually satisfies certain common nonlinearities, such as polynomial, rational, sinusoidal and exponential etc. For instance, diodes are usually modeled by an exponential nonlinear function $i_D = I_s(e^{v_D/(nV_T)} - 1)$ and typical MOSFET models include recursive polynomial, rational and exponential nonlinearities. It is proven in [16] that if the DAE system (1) complies such condition, it can always be converted to another equivalent DAE system with only *polynomial nonlinearities*

$$F_p(x, \dot{x}, \ddot{x}, \dots, u) = \mathbf{0}, \quad (2)$$

where $x \in \mathbb{R}^n$, $u \in \mathbb{R}^l$ and $F_p(\cdot)$ is a polynomial vector function of u, x, \dot{x}, \ddot{x} , etc.

The best way to illustrate the process is by means of an example. Suppose we have a ‘‘diode-type’’ exponential nonlinear system

$$F(z, u) = -\dot{z} - z - e^{4z} + 1 + u = 0, \quad (3)$$

wherein u is a scalar input and z the scalar state variable. By setting $z_1 = z$, $z_2 = e^{4z_1} - 1$ and $x = (z_1, z_2)^T$ such that $\dot{z}_2 = 4e^{4z_1}\dot{z}_1 = 4(z_2 + 1)(-z_1 - z_2 + u)$, we end up with the polynomial DAE system

$$\begin{aligned} F_p(x, \dot{x}, u) &= C_0\dot{x} + G_1x + G_2x \otimes x + D_1x \otimes u + Bu = \mathbf{0} \\ \text{with } C_0 &= \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}, G_1 = \begin{pmatrix} -1 & -1 \\ -4 & -4 \end{pmatrix}, \\ G_2 &= \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & -4 & 0 & -4 \end{pmatrix}, D_1 = \begin{pmatrix} 0 & 0 \\ 0 & 4 \end{pmatrix}, B = \begin{pmatrix} 1 \\ 4 \end{pmatrix}, \end{aligned}$$

where the \otimes symbol denotes the Kronecker product.

For notational ease, we will illustrate our proposed algorithm on the following simplified $F_p(\cdot)^1$

$$\begin{aligned} F_p(x, \dot{x}, \ddot{x}, u) &= C_0\dot{x} + C_1x \otimes \dot{x} + C_2x^{\otimes 2} \otimes \dot{x} + E_0\ddot{x} + G_1x \\ &\quad + G_2x^{\otimes 2} + G_3x^{\otimes 3} + D_1x \otimes u + D_2x^{\otimes 2} \otimes u \\ &\quad + Bu = \mathbf{0}, \end{aligned} \quad (4)$$

where $C_i \in \mathbb{R}^{n \times n^{i+1}}$, $E_0 \in \mathbb{R}^{n \times n}$, $G_i \in \mathbb{R}^{n \times n^i}$, $D_i \in \mathbb{R}^{n \times ln^i}$ and $B \in \mathbb{R}^{n \times l}$ are all constant matrices. All derivations in this paper can easily accommodate higher order powers or cross products of u, x, \dot{x}, \ddot{x} , etc. A formal proof and procedure on how to convert the original DAE system into an equivalent polynomial system can be found in [16].

III. LINEAR ALGEBRA PSS FINDER

A. Harmonic balancing equation

We are now ready to set up a system of multivariate polynomials from (4) by the harmonic balancing method.

Suppose the PSS solution of x in (4) can be approximated by a truncated Fourier series

$$x(t) = \sum_{k=-m}^m p_k e^{jk\omega t}, \quad (5)$$

where $p_k \in \mathbb{C}^n$ are vectors of complex coefficients and ω is a constant if the frequency of the input u is known. For

autonomous systems, ω can be regarded as an extra unknown variable. Since for any real signal $x(t)$, its coefficients in the Fourier bases $e^{jk\omega t}$ and $e^{-jk\omega t}$ must be a pair of complex conjugates, we can specifically define

$$p_0 = p_{0r}, p_{\pm k} = p_{kr} \pm jp_{ki}, k = 1, \dots, m, \quad (6)$$

where $p_{0r}, p_{kr}, p_{ki} \in \mathbb{R}^n$. We further define the coefficients matrix $P = (p_{-m}, p_{-m+1}, \dots, p_0, \dots, p_{m-1}, p_m)$ and the Fourier basis $\Phi_m(t) = (e^{-jm\omega t}, \dots, e^{jm\omega t})^T$, thus $x(t) = P\Phi_m$. Similarly, a periodic input u can always be expressed under the Fourier basis as $u = U\Phi_m$.

The differential operational matrix under the Fourier basis is then defined as

$$D_m = j\omega \text{diag}([-m, -m+1, \dots, m-1, m]). \quad (7)$$

It can be found that the time derivative of the basis functions Φ_m is $D_m\Phi_m$. Therefore \dot{x} and \ddot{x} are simply $PD_m\Phi_m$ and $PD_m^2\Phi_m$, respectively.

The Kronecker power $x^{\otimes 2}$ can be written as

$$x^{\otimes 2} = (P\Phi_m)^{\otimes 2} = (P\Phi_m) \otimes (P\Phi_m) = P^{\otimes 2}\Phi_m^{\otimes 2}. \quad (8)$$

The Kronecker product of the Fourier basis $\Phi_m^{\otimes 2}$ can be simplified to a selector matrix, denoted $S_{m,m}$, times Φ_{2m} . For example the selector matrix $S_{m,m}$ when $m = 1$ is

$$\begin{aligned} \Phi_1^{\otimes 2} &= \begin{pmatrix} e^{-j\omega t} \\ 1 \\ e^{j\omega t} \end{pmatrix} \otimes \begin{pmatrix} e^{-j\omega t} \\ 1 \\ e^{j\omega t} \end{pmatrix} \\ &= \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}^T \begin{pmatrix} e^{-j2\omega t} \\ e^{-j\omega t} \\ 1 \\ e^{j\omega t} \\ e^{j2\omega t} \end{pmatrix} = S_{1,1}\Phi_2. \end{aligned}$$

Similarly, the Kronecker product of the Fourier basis $\Phi_m^{\otimes 3}$ can be simplified to

$$\Phi_m^{\otimes 3} = S_{m,m}\Phi_{2m} \otimes \Phi_m = S_{m,m}S_{2m,m}\Phi_{3m} \triangleq S_{m,m,m}\Phi_{3m}. \quad (9)$$

High order Kronecker products can be constructed in the same way. It should be noticed that instead of growing exponentially, the basis dimension increases only linearly by taking the advantage of selector matrices. Therefore, we have the following equalities

$$\begin{aligned} x &= P\Phi_m, \\ \dot{x} &= PD_m\Phi_m, \\ \ddot{x} &= PD_m^2\Phi_m, \\ x^{\otimes 2} &= P^{\otimes 2}S_{m,m}\Phi_{2m}, \\ x \otimes \dot{x} &= P^{\otimes 2}(I_m \otimes D_m)S_{m,m}\Phi_{2m}, \\ x \otimes u &= (P \otimes U)S_{m,m}\Phi_{2m}, \\ x^{\otimes 3} &= P^{\otimes 3}S_{m,m,m}\Phi_{3m}, \\ x^{\otimes 2} \otimes \dot{x} &= P^{\otimes 3}(I_{2m} \otimes D_m)S_{m,m,m}\Phi_{3m}, \\ x^{\otimes 2} \otimes u &= (P^{\otimes 2} \otimes U)S_{m,m,m}\Phi_{3m}, \\ &\dots \end{aligned}$$

¹We use the shorthand $x \otimes x = x^{\otimes 2}$ etc. throughout the paper.

Replacing the corresponding terms in (4), the original equation can be written in the Fourier basis as

$$\begin{aligned} \mathbf{0} = & C_0 P D_m \Phi_m + C_1 P^{\odot} (I_m \otimes D_m) S_{m,m} \Phi_{2m} \\ & + C_2 P^{\odot} (I_{2m} \otimes D_m) S_{m,m,m} \Phi_{3m} + E_0 P D_m^2 \Phi_m \\ & + G_1 P \Phi_m + G_2 P^{\odot} S_{m,m} \Phi_{2m} + G_3 P^{\odot} S_{m,m,m} \Phi_{3m} \\ & + D_1 (P \otimes U) S_{m,m} \Phi_{2m} + D_2 (P^{\odot} \otimes U) S_{m,m,m} \Phi_{3m} + B U \Phi_m. \end{aligned} \quad (10)$$

Since we are interested in the coefficients with respect to the basis Φ_m , it is necessary to truncate all other bases with an order higher than m or lower than $-m$. The truncation can be achieved by removing the columns in each S that correspond to elements that are not in the basis Φ_m . Suppose the new selector matrix is \tilde{S} after the truncation, (10) then becomes

$$\begin{aligned} \mathbf{0} = & C_0 P D_m \Phi_m + C_1 P^{\odot} (I_m \otimes D_m) \tilde{S}_{m,m} \Phi_m \\ & + C_2 P^{\odot} (I_{2m} \otimes D_m) \tilde{S}_{m,m,m} \Phi_m + E_0 P D_m^2 \Phi_m \\ & + G_1 P \Phi_m + G_2 P^{\odot} \tilde{S}_{m,m} \Phi_m + G_3 P^{\odot} \tilde{S}_{m,m,m} \Phi_m \\ & + D_1 (P \otimes U) \tilde{S}_{m,m} \Phi_m + D_2 (P^{\odot} \otimes U) \tilde{S}_{m,m,m} \Phi_m + B U \Phi_m. \end{aligned} \quad (11)$$

The Φ_m is a common factor to each term in (11) and can hence be removed to arrive at the harmonic balancing equation

$$\begin{aligned} \mathbf{0} = & C_0 P D_m + C_1 P^{\odot} (I_m \otimes D_m) \tilde{S}_{m,m} \\ & + C_2 P^{\odot} (I_{2m} \otimes D_m) \tilde{S}_{m,m,m} + E_0 P D_m^2 \\ & + G_1 P + G_2 P^{\odot} \tilde{S}_{m,m} + G_3 P^{\odot} \tilde{S}_{m,m,m} \\ & + D_1 (P \otimes U) \tilde{S}_{m,m} + D_2 (P^{\odot} \otimes U) \tilde{S}_{m,m,m} + B U. \end{aligned} \quad (12)$$

It should be noted that both sides of (12) are n -by- $(2m+1)$ matrices. Therefore, after separating the real and imaginary parts of (12), (12) can be regarded as a system of $n(2m+1)$ multivariate polynomials with $n(2m+1)$ unknowns (p_{0r} , p_{kr} and p_{ki}).

For autonomous cases, the extra variable ω needs to be introduced, which makes (12) underdetermined. This results in an infinite amount of solutions, which come from two sources. First, for autonomous systems, any arbitrary phase shift of a solution to (12) will also be a valid solution to the original system. Therefore, to enforce the solutions be isolated from one another, it is necessary to fix the phase of the solution by, e.g., setting the first element² of p_{1i} to 0 (forcing the phase of the first order harmonic of the first state variable to be 0 or π). Note that setting p_{1i} to 0 in (12) and the addition of ω results in a polynomial system of $n(2m+1)$ equations and unknowns. Second, it is readily found that each harmonic order of the fundamental frequency is also a valid oscillating frequency. In other words, if $(\omega, p_1, p_2, p_3, p_4, \dots)$ is a set of solutions to (12), then so are $(2\omega, 0, p_1, 0, p_2, 0, p_3, \dots)$ and $(3\omega, 0, 0, p_1, 0, 0, p_2, \dots)$, etc. Thus, to ensure ω to be the fundamental oscillating frequency, we set the first element of p_{1r} nonzero by adding an extra polynomial equation $h \cdot p_{1r}^1 - 1 = 0$, where h is an extra variable.

In this way autonomous systems will have $n(2m+1)+1$ polynomials and unknowns, while forced systems will have $n(2m+1)$ polynomials and unknowns. For the remainder of this article we denote the number of polynomials and unknowns by \hat{n} and consider systems of \hat{n} polynomials in the variables $x_1, \dots, x_{\hat{n}}$.

The procedure of building the harmonic balancing polynomial system is summarized in Algorithm 1.

Algorithm 1 Polynomial system builder

Input: $F(x, \dot{x}, \ddot{x}, \dots, u)$, n , m

Output: system of \hat{n} polynomials in \hat{n} unknowns

- 1: convert $F(x, \dot{x}, \ddot{x}, \dots, u) = 0$ into an equivalent polynomial system $F_p(x, \dot{x}, \ddot{x}, \dots, u) = 0$
 - 2: get B, C_i, E_i, G_i, D_i , etc. from F_p
 - 3: build an n -by- $2m+1$ coefficient matrix P
 - 4: apply P, B, C_i, E_i, G_i, D_i , etc. in (12)
 - 5: get the multivariate polynomial system $f_1, \dots, f_{n(2m+1)}$
 - 6: **if** autonomous **then**
 - 7: $\hat{n} = n(2m+1) + 1$
 - 8: $p_{1i}^1 = 0$
 - 9: add an equation $f_{\hat{n}}: h p_{1r}^1 - 1 = 0$
 - 10: **else**
 - 11: $\hat{n} = n(2m+1)$
 - 12: **end if**
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B. Multivariate polynomial root-finder

We have established in the previous section that the coefficients of a truncated Fourier series of the periodic state solution can be computed as the roots of a multivariate polynomial system. The next step now is to find all the roots of this multivariate polynomial system. The main ingredient of our numerical root finder, which is partly described in [13], [14], is the Macaulay matrix $M(d)$. For a given degree $d \in \mathbb{N}$ and system of \hat{n} -variate polynomials $f_1, \dots, f_{\hat{n}}$ of degrees $d_1, \dots, d_{\hat{n}}$, the Macaulay matrix $M(d)$ contains the coefficients of

$$\begin{pmatrix} f_1 \\ x_1 f_1 \\ \vdots \\ x_{\hat{n}}^{d-d_1} f_1 \\ f_2 \\ x_1 f_2 \\ \vdots \\ x_{\hat{n}}^{d-d_s} f_{\hat{n}} \end{pmatrix}, \quad (13)$$

where each polynomial f_i is multiplied with all monomials from degree 0 up to $d - d_i$ for all $i = 1, \dots, \hat{n}$. Each of these polynomials $f_1, \dots, f_{\hat{n}}$ can be made homogeneous by adding an additional variable x_0 such that each term has the same degree. The roots of the homogeneous polynomials are then described by the projective coordinates $(x_0, x_1, \dots, x_{\hat{n}})$. Roots at infinity are roots for which $x_0 = 0$ and affine roots are roots for which $x_0 \neq 0$. We will denote the homogenization of f_1 by f_1^h and $\mathcal{P}_d^{\hat{n}}$ will denote the vector space of \hat{n} -variate homogeneous polynomials of degree d . The row space of $M(d)$ is then the vector space

$$\mathcal{M}_d = \left\{ \sum_{i=1}^{\hat{n}} h_i f_i^h \mid h_i \in \mathcal{P}_{d-d_i}^{\hat{n}} \right\}. \quad (14)$$

Let $q(d), r(d), c(d)$ respectively denote the number of columns, the rank and the nullity of $M(d)$. Then the rank-nullity theorem can be written as

$$q(d) = r(d) + c(d), \quad (15)$$

²We denote p_{1i}^1 to be the first element of the vector p_{1i} , etc.

where $q(d) = \binom{d+\hat{n}}{\hat{n}} = \dim \mathcal{P}_d^{\hat{n}}$ and $r(d) = \dim \mathcal{M}_d$. By rewriting (15), the following expression for $c(d)$ is found

$$\begin{aligned} c(d) &= q(d) - r(d), \\ &= \dim \mathcal{P}_d^n - \dim \mathcal{M}_d, \\ &= \dim \mathcal{P}_d^n / \mathcal{M}_d. \end{aligned}$$

The function in the degree d that describes the dimension of the quotient space $\mathcal{P}_d^n / \mathcal{M}_d$ is called the Hilbert Function [17, p. 462]. It can be shown that there is a degree d^* for which the Hilbert Function becomes a polynomial and its degree corresponds with the dimension of the projective solution set [17, p. 464]. This implies that for a polynomial system with a finite number of s projective solutions we have that $c(d) = s$ ($\forall d \geq d^*$). It is also possible that the number of affine roots is finite while those at infinity are not finite in number. In this case $c(d)$ will grow polynomially from a certain degree d^* . We assume from here on, without loss of generality, that there are a finite number of projective roots and that none of the s roots z_1, \dots, z_s have any multiplicities. The canonical null space K of $M(d)$ is then

$$K = (\delta_0|_{z_1} \quad \delta_0|_{z_2} \quad \dots \quad \delta_0|_{z_s})$$

with

$$\delta_0|_z = (x_0^d \quad x_0^{d-1}x_1 \quad x_0^{d-1}x_2 \quad \dots \quad x_{\hat{n}}^d)^T$$

in the root $z = (x_0, x_1, \dots, x_{\hat{n}})$. Note that for an affine root, one can set $x_0 = 1$ in $\delta_0|_z$ and obtain

$$\delta_0|_z = (1 \quad x_1 \quad x_2 \quad \dots \quad x_{\hat{n}}^d)^T.$$

This affine canonical null space vector has an \hat{n} -variate Vandermonde structure, which implies the following shift property

$$(1 \quad x_1 \quad x_2 \quad \dots \quad x_{\hat{n}}^d)^T x_1 = (x_1 \quad x_1^2 \quad x_1x_2 \quad \dots \quad x_1x_{\hat{n}}^d)^T. \quad (16)$$

Evidently, this shift property can be written for multiplication with any monomial x_i . In addition, (16) can be expressed completely in terms of $\delta_0|_z$ by limiting it to the rows of $\delta_0|_z$ corresponding with monomials of degree at most $d-1$. Let S_1 be the matrix that selects the aforementioned rows and S_{x_i} the matrix that selects the corresponding rows after multiplication with x_i . Then (16) can be written as

$$S_1 \delta_0|_z x_i = S_{x_i} \delta_0|_z.$$

The canonical kernel K can be partitioned into

$$K = \begin{pmatrix} s_a & s - s_a \\ K_a & K_\infty \end{pmatrix}$$

where K_a contains the $\delta_0|_z$ vectors for the affine roots and K_∞ those for the roots at infinity. The shift property can then be written for all s_a affine roots as

$$S_1 K_a D_{x_i} = S_{x_i} K_a, \quad (17)$$

with D_{x_i} a $s_a \times s_a$ diagonal matrix containing the x_i component of the corresponding affine roots. The K_a matrix is not known however. Instead, a numerical basis N for the null space can be computed using either the singular value decomposition (SVD) or rank-revealing QR [15]. The SVD is the more robust choice for the determination of the rank of $M(d)$ but

computationally more expensive. Obviously, $K = NT$ and therefore $K_a = NT_a$. We now partition N into

$$\begin{pmatrix} s \\ k & \begin{pmatrix} N_1 \\ N_2 \end{pmatrix} \\ q - k \end{pmatrix}$$

such that $\text{rank}(N_1) = s_a$ and compute orthogonal bases Q_1, Q_2 such that $\text{col}(Q_1) = \text{col}(N_1)$ and $\text{col}(Q_2) = \text{null}(N_1)$. This can also be done using either an SVD or rank-revealing QR. Then, the following basis change can be performed $Z = N \begin{pmatrix} Q_1 & Q_2 \end{pmatrix}$. This can be written out as

$$Z = N Q = \begin{pmatrix} k & \begin{pmatrix} N_1 \\ N_2 \end{pmatrix} \\ q - k \end{pmatrix} \begin{pmatrix} s_a & s - s_a \\ Q_1 & Q_2 \end{pmatrix}.$$

Since Q_2 is orthogonal to N_1 we have that

$$Z = \begin{pmatrix} k & \begin{pmatrix} N_1 Q_1 & 0 \\ N_2 Q_1 & N_2 Q_2 \end{pmatrix} \\ q - k \end{pmatrix} \begin{pmatrix} s_a & s - s_a \\ Q_1 & Q_2 \end{pmatrix}. \quad (18)$$

Observe that the upper right block of Z is zero, for this reason this linear basis change is also called a column compression. Also, since $K = NT$ and $Z = NQ$, we can write $K = ZV$ with $V = Q^T T$. By applying the same partitioning as (18), $K = ZV$ can be rewritten as

$$\begin{pmatrix} K_{a1} & K_{\infty 1} \\ K_{a2} & K_{\infty 2} \end{pmatrix} = \begin{pmatrix} k & \begin{pmatrix} N_1 Q_1 & 0 \\ N_2 Q_1 & N_2 Q_2 \end{pmatrix} \\ q - k \end{pmatrix} \begin{pmatrix} s_a & s - s_a \\ V_{11} & V_{12} \\ V_{21} & V_{22} \end{pmatrix},$$

from which we find that $K_{a1} = N_1 Q_1 V_{11}$. Hence, by choosing S_1 and S_{x_i} in (17) such that only rows of K_{a1} are selected we can write (17) as

$$S_1 K_{a1} D_{x_i} = S_{x_i} K_{a1}.$$

Substituting K_{a1} this becomes

$$S_1 N_1 Q_1 V_{11} D_{x_i} = S_{x_i} N_1 Q_1 V_{11}.$$

Setting $B \triangleq S_1 N_1 Q_1$ and $A \triangleq S_{x_i} N_1 Q_1$, we obtain

$$B V_{11} D_{x_i} = A V_{11}, \quad (19)$$

which is a rectangular eigenvalue problem. This can be solved as

$$V_{11} D_{x_i} V_{11}^{-1} = B^\dagger A$$

where B^\dagger is the Moore-Penrose pseudoinverse of B . Alternatively, one could choose S_1, S_{x_i} also such that B, A are square and B of full rank. Then (19) is a generalized eigenvalue problem. Once the eigenvectors V_{11} are computed, they can be used to reconstruct the affine part of the canonical null space K_a since $K_{a1} = N_1 Q_1 V_{11}$. All the affine roots can then be read off after scaling the first row of the obtained K_{a1} such that all its elements are 1. We summarize the pseudo code for the proposed root-finding method in Algorithm 2.

As described in [13], the number of affine roots s_a can be determined from monitoring the linearly independent leading monomials in \mathcal{M}_d . Indeed, a sufficient condition for the polynomial system to have a finite number of affine roots is that for each variable x_i ($i = 1, \dots, \hat{n}$), some pure power x_i^α is present in the linearly independent monomials of \mathcal{M}_d . The computational complexity of Algorithm 2 is dominated

by the computation of a basis for the null space of $M(d)$. For the SVD or rank-revealing QR this complexity is $O(pq^2)$ where p is the number of rows and q the number of columns of $M(d)$. In terms of \hat{n} this complexity is $O(\hat{n}d^{3\hat{n}}/(\hat{n}!)^3)$. Using a graded monomial ordering for the columns of $M(d)$ results in a particular matrix structure that can be exploited to reduce the complexity to $O((\hat{n}-1)d^{3\hat{n}-3}/(\hat{n}-1!)^3)$ [18]. This optimization however does not take the quasi-Toeplitz structure nor the sparsity of the Macaulay matrix into account. Indeed, the density of $M(d)$, which is the ratio of the number of nonzero elements with the total number of entries, decreases as $O((\hat{n}-1)!/(d^{\hat{n}}))$. Further optimizations will hence involve the development of either direct or iterative methods to compute the null space of the sparse and structured Macaulay matrix.

Algorithm 2 Affine root-finder

Input: \hat{n} -variate polynomials $f_1, \dots, f_{\hat{n}}$

Output: all affine roots of $f_1, \dots, f_{\hat{n}}$

- 1: construct $M(d)$ until $c(d)$ stabilizes and all linearly independent pure powers x_i^α are found
 - 2: compute numerical basis N for null space of $M(d)$
 - 3: perform column compression (18) on N
 - 4: choose S_1, S_{x_i} and solve eigenvalue problem (19)
 - 5: reconstruct affine canonical kernel by $K_{a1} = N_1 Q_1 V_{11}$
 - 6: scale K_{a1} and read off all affine roots
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IV. NUMERICAL EXAMPLES

In this section we illustrate our proposed scheme and compare it with HB by means of several numerical examples. All algorithms were implemented in MATLAB and done on a desktop PC @ 3.0 GHz and 16 GB RAM. In addition, we also compare our computed roots of the polynomial systems with those computed in Maple using Groebner bases.

A. van der Pol oscillator

Let us consider the following van der Pol-like equation

$$F_p(x) = \ddot{x} - (1 - 3x^2)\dot{x} + x = 0. \quad (20)$$

An electrical circuit implementation for this equation is shown in Figure 1. The tunnel diode has the cubic current-voltage relationship

$$i = \phi_t(v) = \phi(v - E_0) + I,$$

with $\phi(v) = \gamma v^3 - \alpha v$. The equation for the circuit in Figure 1 can then be written as

$$\dot{V} = \frac{1}{C}(-\phi(V) - W)$$

with $\dot{W} = V/L$. This can be rewritten to

$$\ddot{V} - \frac{1}{C}(\alpha - 3\gamma V^2)\dot{V} + \frac{1}{LC}V = 0. \quad (21)$$

Setting $\alpha = \gamma = 1$ and $C = 1F, L = 1H$ results in (20).

We set the order of the Fourier approximation to $m = 3$ and use the harmonic balancing scheme from Subsection III-A. Due to the odd symmetry property of (20), all even order

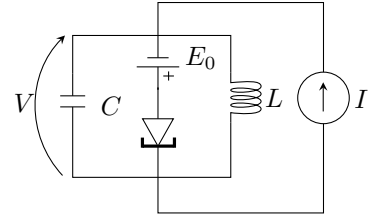


Fig. 1. An electrical circuit with a tunnel diode for the van der Pol oscillator.

Fourier coefficients of x must be 0 (i.e., $p_0 = p_{\pm 2} = 0$). This results in the following polynomial system

$$\begin{cases} -\omega^2 - 3p_{1r}p_{3i}\omega + 1 = 0, \\ 6p_{3r}^2 + 3p_{1r}p_{3r} + 6p_{3i}^2 + 3p_{1r}^2 - 1 = 0, \\ -9p_{3r}\omega^2 - 9p_{3i}p_{3r}^2\omega - 9p_{3i}^3\omega - 18p_{1r}^2p_{3i}\omega \\ + 3p_{3i}\omega + p_{3r} = 0, \\ 9p_{3r}^3\omega - 9p_{3i}\omega^2 + 9p_{3i}^2p_{3r}\omega + 18p_{1r}^2p_{3r}\omega \\ - 3p_{3r}\omega + 3p_{1r}^3\omega + p_{3i} = 0, \end{cases}$$

of 4 polynomials in 4 unknowns $(\omega, p_{1r}, p_{3r}, p_{3i})$. For this system $c(d)$ grows linearly from $d \geq 7$. Only at $d = 9$ all pure powers are present in the linearly independent leading monomials. From the linearly independent leading monomials we deduce that the number of affine roots $s_a = 20$. Note that the linear increase of $c(d)$ together with $s_a = 20$ implies that solutions set of the roots at infinity is 1-dimensional. $M(9)$ is a 792×715 matrix with a rank of 607 and a density of 0.7%. This means that $M(9)$ consists of more than 99% zeros. The computed numerical basis for the null space N is 715×108 . If we choose S_1 such that it selects all rows of N corresponding with monomials of degree at most 3, then we have that $\text{rank}(N_1) = 20$. Performing the column compression, solving the eigenvalue problem and reconstructing the kernel results in 20 affine roots $(\omega, p_{1r}, p_{3r}, p_{3i})$. Only 4 have a real ω , they are all the same solution within a sign change

$$\begin{pmatrix} 0.9426, & 0.5825, & -0.0288, & 0.0676 \\ -0.9426, & -0.5825, & 0.0288, & 0.0676 \\ 0.9426, & -0.5825, & 0.0288, & -0.0676 \\ -0.9426, & 0.5825, & -0.0288, & -0.0676 \end{pmatrix}.$$

The 2-norm of the vector containing the residuals is of order of magnitude 10^{-13} . The absolute difference between our computed roots and the roots from the Groebner basis is also of order of magnitude 10^{-13} .

Next, we apply HB on the oscillator. If the initial guess is carefully chosen to be sufficiently close to the desired solution, e.g. $(1, 0.5, 0, 0)$, then HB converges to the PSS solution in 9 iterations if a tolerance of 10^{-13} is used. However, if an unfortunate “bad” initial guess such as $(0.2, 1, 0, 0)$ or $(0.94, 0.2, 0, 0)$ is used, HB will diverge after a few iterations, even when ω is close to the solution in $(0.94, 0.2, 0, 0)$.

B. Wien bridge oscillator

The second example is a Wien bridge oscillator adapted from [10]. The circuit and its controlled source model are

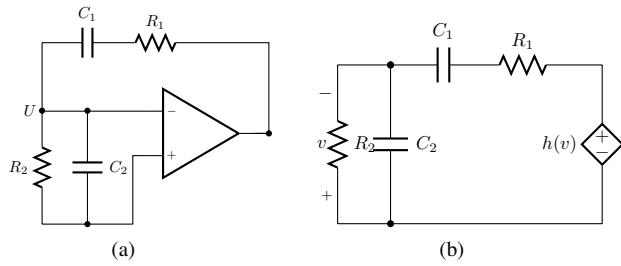


Fig. 2. (a) A Wien bridge oscillator circuit. (b) Controlled source model of the oscillator.

shown in Figure 2, where $R_1 = R_2 = 1\Omega$ and $C_1 = C_2 = 1F$. The voltage-controlled voltage source is described by

$$h(v) = 3.234v - 2.195v^3 + 0.666v^5.$$

Its univariate DAE can be obtained

$$F_p(v) = \ddot{v} - 0.234\dot{v} + v + 6.585v^2\dot{v} - 3.334v^4\dot{v} = 0. \quad (22)$$

Similarly, (22) is odd symmetric, therefore all even order Fourier coefficients should also be 0. Applying the harmonic balancing scheme for $m = 3$ results in the polynomial system shown at the top of the next page.

The dimension of the null space $c(d)$ also grows linearly from $d \geq 14$. Only at $d = 16$ all pure powers are present in the linearly independent leading monomials. Hence the number of affine roots is again finite and determined to be $s_a = 80$. $M(16)$ is a 5187×4845 matrix with a rank of 3999 and a density of 0.2%. The computed numerical basis for the null space N is 4845×846 . Choosing S_1 such that it selects all rows of N corresponding with monomials of degree at most 5 ensures that $\text{rank}(N_1) = 80$. Performing the column compression, solving the eigenvalue problem and reconstructing the kernel results in 80 affine roots $(\omega, p_{1r}, p_{3r}, p_{3i})$. Only 8 have a real ω , they are

$$\begin{pmatrix} 0.7976, & -1.4856, & 0.4119, & -0.0191 \\ 0.7976, & 1.4856, & -0.4119, & -0.0191 \\ -0.7976, & 1.4856, & -0.4119, & 0.0191 \\ 0.7976, & -1.4856, & 0.4119, & 0.0191 \\ 0.9967, & 0.1922, & -0.0005, & 0.0055 \\ -0.9967, & -0.1922, & 0.0005, & 0.0055 \\ 0.9967, & -0.1922, & 0.0005, & -0.0055 \\ -0.9967, & 0.1922, & -0.0005, & -0.0055 \end{pmatrix}$$

The 2-norm of the vector containing the residuals is of order of magnitude 10^{-13} . The absolute difference between our computed roots and the roots from the Groebner basis is of order of magnitude 10^{-12} . Observe that these results suggest that 2 PSSs exist: one with an angular frequency of 0.7976 and one with an angular frequency of 0.9967. Simulations show that the PSS with $\omega = 0.7976$ is unstable, while the other PSS is stable.

Using HB, the stable PSS can be obtained if a “good” initial value, e.g., $(1, 0.2, 0, 0)$, is given. Otherwise, for instance using the initial guess $(1, 1, 0, 0)$ HB diverges. Furthermore, if $(0.2, 1, 0, 0)$ is chosen, HB will converge to the unstable solution with $\omega = 0.7976$.

C. Biochemical reaction system

Then we illustrate a multivariate example with a rational nonlinearity. The system is used to describe a biochemical reaction adapted from [16]. It is an excited system and can be expressed by

$$\dot{x} + x + \frac{x}{1 + 10x} + u = 0, \quad (23)$$

where $u = \cos(2\pi t)$ is the external input. Using a Taylor expansion is inaccurate to approximate the rational nonlinearity since x is not guaranteed to be within the convergent range of the Taylor expansion. In our scheme, by adding the extra variable $y = x/(1 + 10x)$ such that $-x + y + 10xy = 0$, we come up with the polynomial DAE system $F_p(\dot{x}, x, y, u) = 0$

$$\begin{cases} \dot{x} + x + y + u = 0, \\ -x + y + 10xy = 0. \end{cases} \quad (24)$$

Using $m = 2$ in the Fourier approximation results in a polynomial system of 10 equations in 10 unknowns. We have omitted these polynomials due to lack of space. At $d = 7$ all pure powers are present in the linearly independent leading monomials. The number of affine roots is determined to be $s_a = 12$. $M(7)$ is a 55055×19448 matrix with a rank of 19425 and a density of 0.02%. Choosing S_1 such that it selects all rows of N corresponding with monomials of degree at most 2 ensures that $\text{rank}(N_1) = 12$. Performing the column compression, solving the eigenvalue problem and reconstructing the kernel results in 12 affine roots, of which 4 are real. They are given by

$$\begin{pmatrix} p_{0r}^x \\ p_{0r}^y \\ p_{1r}^x \\ p_{1r}^y \\ p_{2r}^x \\ p_{2r}^y \\ p_{1i}^x \\ p_{1i}^y \\ p_{2i}^x \\ p_{2i}^y \end{pmatrix} = \left\{ \begin{pmatrix} 0.0551 \\ -0.0551 \\ -0.0251 \\ -0.0368 \\ -0.0028 \\ 0.0307 \\ 0.0697 \\ 0.0878 \\ 0.0022 \\ 0.0330 \end{pmatrix}, \begin{pmatrix} -0.1485 \\ 0.1485 \\ -0.0049 \\ 0.0006 \\ -0.0014 \\ 0.0765 \\ 0.0789 \\ -0.0479 \\ 0.0060 \\ 0.0118 \end{pmatrix}, \begin{pmatrix} -0.2551 \\ 0.2551 \\ -0.0251 \\ -0.0368 \\ 0.0028 \\ -0.0307 \\ 0.0697 \\ 0.0878 \\ -0.0022 \\ -0.0330 \end{pmatrix}, \begin{pmatrix} -0.0515 \\ 0.0515 \\ -0.0050 \\ 0.0006 \\ 0.0014 \\ -0.0765 \\ 0.0789 \\ -0.0479 \\ -0.0060 \\ -0.0118 \end{pmatrix} \right\}.$$

It can be shown that only the first two columns of solutions satisfy the hypothesis $y = x/(1 + 10x)$. Further transient simulation shows that the first column is the PSS solution to the forced system. The 2-norm of the vector containing the residuals is of order of magnitude 10^{-11} . The absolute difference between our computed roots and the roots from the Groebner basis is of order of magnitude 10^{-10} .

Similarly, different initial guesses lead to different results in HB. This is particularly problematic for this example with the occurrence of 4 solutions of which only 1 is desired. For example, $(1, \frac{1}{11}, 0, \dots, 0)$ can drive to the solution in the second column while $(0, \dots, 0)$ results in the right PSS solution.

D. Colpitts oscillator

Finally we study an NMOS Colpitts oscillator shown in Figure 3. It is a typical example for autonomous circuits. We use the same normalized state variables as defined in [19]. The normalized state equations of the Colpitts oscillator can

$$\begin{cases} 3.334 p_{3i}^3 p_{1r} \omega + 10.002 p_{3i} p_{1r}^3 \omega + 3.334 p_{3i} p_{1r} p_{3r}^2 \omega - 6.585 p_{3i} p_{1r} \omega - 1.0 \omega^2 + 1.0 = 0, \\ 13.336 p_{3i}^4 + 6.668 p_{3i}^2 p_{1r}^2 + 3.334 p_{3i}^2 p_{1r} p_{3r} + 26.672 p_{3i}^2 p_{3r}^2 - 13.17 p_{3i}^2 + 6.668 p_{1r}^4 + 16.67 p_{1r}^3 p_{3r} + 6.668 p_{1r}^2 p_{3r}^2 - 6.585 p_{1r}^2 \\ + 3.334 p_{1r} p_{3r}^3 - 6.585 p_{1r} p_{3r} + 13.336 p_{3r}^4 - 13.17 p_{3r}^2 + 0.234 = 0, \\ 10.002 p_{3i}^5 \omega + 40.008 p_{3i}^3 p_{1r}^2 \omega + 20.004 p_{3i}^3 p_{3r}^2 \omega - 19.755 p_{3i}^3 \omega + 43.342 p_{3i} p_{1r}^4 \omega + 6.668 p_{3i} p_{1r}^3 p_{3r} \omega + 40.008 p_{3i} p_{1r}^2 p_{3r}^2 \omega \\ - 39.51 p_{3i} p_{1r}^2 \omega + 10.002 p_{3i} p_{3r}^4 \omega - 19.755 p_{3i} p_{3r}^2 \omega + 0.702 p_{3i} \omega - 9.0 p_{3r} \omega^2 + p_{3r} = 0, \\ 10.002 p_{3i}^4 p_{3r} \omega + 3.334 p_{3i}^2 p_{1r}^3 \omega + 40.008 p_{3i}^2 p_{1r}^2 p_{3r} \omega + 20.004 p_{3i}^2 p_{3r}^3 \omega - 19.755 p_{3i}^2 p_{3r} \omega + 9.0 p_{3i} \omega^2 - 1.0 p_{3i} + 43.342 p_{1r}^4 p_{3r} \omega \\ + 10.002 p_{1r}^5 \omega + 10.002 p_{1r}^3 p_{3r}^2 \omega - 6.585 p_{1r}^3 \omega + 40.008 p_{1r}^2 p_{3r}^3 \omega - 39.51 p_{1r}^2 p_{3r} \omega + 10.002 p_{3r}^5 \omega - 19.755 p_{3r}^3 \omega + 0.702 p_{3r} \omega = 0. \end{cases}$$

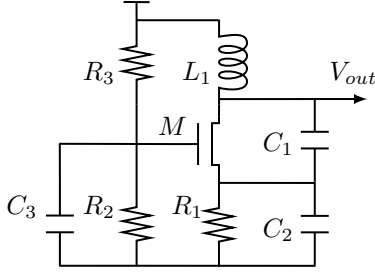


Fig. 3. An NMOS Colpitts oscillator.

be written in the form

$$\begin{cases} \dot{x}_1 = \frac{3}{2}(x_2 + x_3) - \frac{3}{4}x_2^2, \\ \dot{x}_2 = \frac{3}{2}x_3, \\ \dot{x}_3 = -\frac{1}{3}(x_1 + x_2) - \frac{1}{2}x_3. \end{cases}$$

To verify the existence of the periodic solution, we apply our scheme for $m = 1$ and this results in a polynomial system of 10 equations. Here, we also omitted these equations due to the lack of space. It is possible to reduce this system. Indeed, from the polynomial system it is immediately clear that $p_{0r}^3 = 0$ and this also allows us to substitute p_{0r}^2 by $-p_{0r}^1$, effectively reducing the total number of polynomials and unknowns to 8. The dimension of the null space $c(d)$ grows cubically from $d \geq 4$, which implies that the solution set at infinity is 3-dimensional. All 6 affine solutions are found at $d = 8$, at which the Macaulay matrix is 24024×12870 with a density of 0.02%. Each of the 6 affine roots is real. After removing the solutions within sign changes, only the following 2 solutions remain

$$\begin{pmatrix} \omega \\ h \\ p_{0r}^1 \\ p_{0r}^2 \\ p_{0r}^3 \\ p_{1r}^1 \\ p_{1r}^2 \\ p_{1r}^3 \\ p_{1i}^2 \\ p_{1i}^3 \end{pmatrix} = \left\{ \begin{pmatrix} 1.0000 \\ 1.3416 \\ -0.3333 \\ 0.3333 \\ 0.0000 \\ 0.7454 \\ 0.3727 \\ -0.2485 \\ 0.3727 \\ 0.2485 \end{pmatrix}, \begin{pmatrix} 0.0000 \\ 1.4142 \\ -1.0000 \\ 1.0000 \\ 0.0000 \\ 0.7071 \\ -0.7071 \\ 0.0000 \\ 0.0000 \\ 0.0000 \end{pmatrix} \right\}.$$

The 2-norm of residuals is of order of magnitude 10^{-14} and the absolute difference between our computed roots and the roots from the Groebner basis is also of order of magnitude 10^{-14} . The result suggests that there are 2 possible PSS for

the Colpitts oscillator. The first column of the solution is the desired PSS, with $\omega = 1$ satisfying the normalization condition. The second column indicates the unstable constant equilibrium point of the Colpitts oscillator, which can be easily destroyed by a small perturbation.

Not surprisingly, different initial conditions lead to different solutions in HB. For example, $(1, 1, 0, 0, 0, 1, -1, 0, 0, 0)$ leads to the unstable equilibrium point while $(1, 1, 1, 0, 0, 1, 0, 0, 0, 0)$ blows up the iteration. Finally HB converges to the right PSS solution starting from the initial guess $(1, 1, 0, 0, 0, 1, 1, 0, 0, 0)$.

V. REMARKS

Several important remarks are in order:

- 1) Although our balancing equation (12) resembles the harmonic balance (HB) method [3], [4], [7], it is important to observe that HB deals with nonlinear functions by evaluating in the time domain, followed by fast Fourier transforms (FFT) in each iteration of Newton's method. In the proposed method, the nonlinearity is embedded in the polynomial balancing formulation and no domain transforms are required.
- 2) The Groebner basis with respect to the degree negative lex monomial ordering for the Wien bridge example has integer coefficients with up to 150 digits and requires multiprecision integer arithmetic. The proposed method works in double precision and is able to compute results with similar accuracy.
- 3) In subsection II-A, it should be noticed that the size of the state space in (1) will be enlarged after the polynomialization process, i.e., $n > \tilde{n}$ in (1) and (2). The number of extra state variables depends on the number of nonlinearities of each electrical component (including recursive nonlinearities). In practical circuit modeling the number of nonlinearities of each component is always limited and countable. Therefore the complexity of the algorithm will not be increased due to the conversion.
- 4) If the obtained accuracy of our proposed method is not sufficient, then only a few Newton iterations, using the obtained solutions as initial guess, are sufficient to further improve the accuracy up to the machine precision.
- 5) Future work will focus on solving higher order Fourier approximations. This will result in larger polynomial systems, which will be the main bottleneck for the proposed method. Exploiting the structure and nonzero pattern of the Macaulay matrix is then mandatory.

VI. CONCLUSION

In this paper, a novel method for solving the PSS problem of nonlinear oscillators has been proposed. Prior nonlinear autonomous system solvers suffer from either convergence issues or intrinsic coefficient problems, which are circumvented by the proposed method via a Macaulay matrix based polynomial solver. Moreover, exact polynomial DAE systems are constructed from general nonlinear problems to guarantee the strong nonlinearities of the original systems. Several examples confirm the accuracy and feasibility of the proposed method.

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